

# **Karnaugh Maps for Simplification of Digital Logic Functions**

**CS 64: Computer Organization and Design Logic  
Lecture #12  
Winter 2019**

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# Administrative

- Lab #6 on Thursday
  - You don't have to go to lab  
(so this time, we won't take attendance)
  - But the TAs will be there if you need help!
  - Due by **Monday**

# Digital Circuit Design Process

CAN THIS PROCESS BE REVERSED?

Function  
definition

adder



Truth  
table

A	B	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

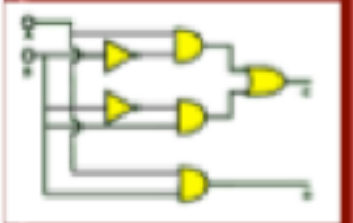


Boolean  
expression

$$\text{Sum} = (A\bar{B}) + (\bar{A}B)$$
$$\text{Carry} = AB$$



Logic  
block



# More Simplification Examples

Simplify the Boolean expression:

- $(A+B+C).(D+E)' + (A+B+C).(D+E)$

Simplify the Boolean expression and write it out on a truth table as proof

- $X.Z + Z.(X' + X.Y)$

Use DeMorgan's Theorem to re-write the expression below using at least one OR operation

- $\text{NOT}(X + Y.Z)$

# Scaling Up Simplification

- When we get to *more* than 3 variables, it becomes challenging to use truth tables
- We can instead use ***Karnaugh Maps*** to make it immediately apparent as to what can be simplified

# Example of a K-Map

	A	B	f(A,B)
0	0	0	a
1	0	1	b
2	1	0	c
3	1	1	d

B \ A	0	1
0	a	c
1	b	d

B \ A	0	1
0	0	2
1	1	3

A	B	f(A,B)
0	0	0
0	1	1
1	0	1
1	1	1

B \ A	0	1
0	0	1
1	1	1

# K-Maps with 3 or 4 Variables

		<i>AB</i>		<i>A</i>	
		00	01	11	10
<i>C</i>	0	0	2	6	4
	1	1	3	7	5

*B*

Note the adjacent placement of:

00 01 11 10

It's NOT:

00 01 10 11

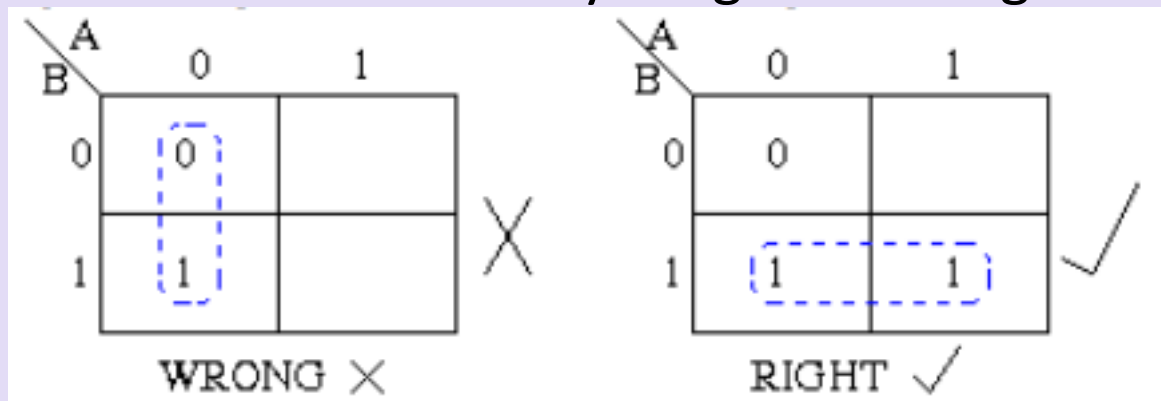
		<i>AB</i>		<i>A</i>	
		00	01	11	10
<i>CD</i>	00	0	4	12	8
	01	1	5	13	9
<i>C</i>	11	3	7	15	11
	10	2	6	14	10

*B*

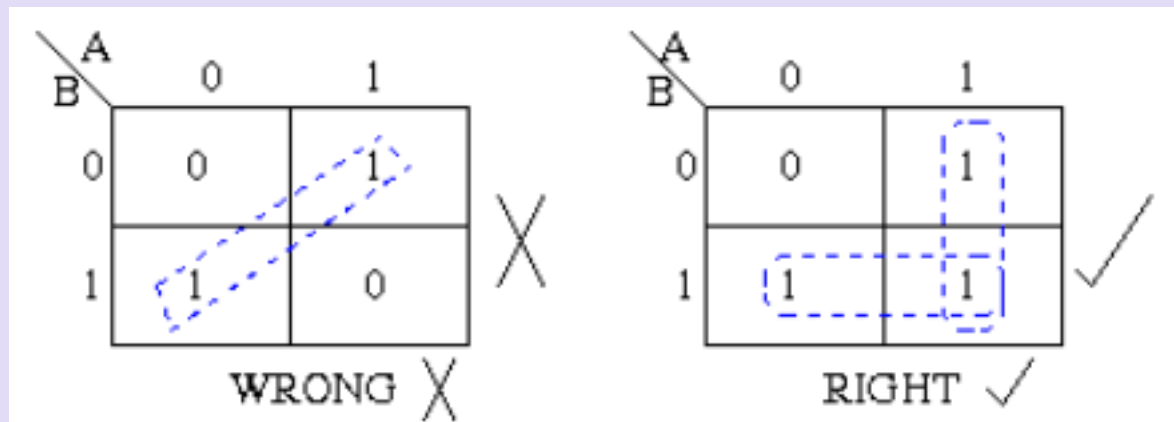
*D*

# Rules for Using K-Maps for Simplification

1. Group together **adjacent cells** containing "1"
2. Groups should **not include** anything containing "0"



3. Groups may be horizontal or vertical, but **not diagonal**





# Rules for Using K-Maps for Simplification

4. Groups must contain 1, 2, 4, 8, or in general  $2^n$  cells.

A B	0	1	
0	1	1	← Group of 2
1	0	0	

RIGHT ✓

AB C	00	01	11	10	
0	0	1	1	1	← Group of 3
1	0	0	0	0	

WRONG ✗

A B	0	1	
0	1	1	← Group of 4
1	1	1	

RIGHT ✓

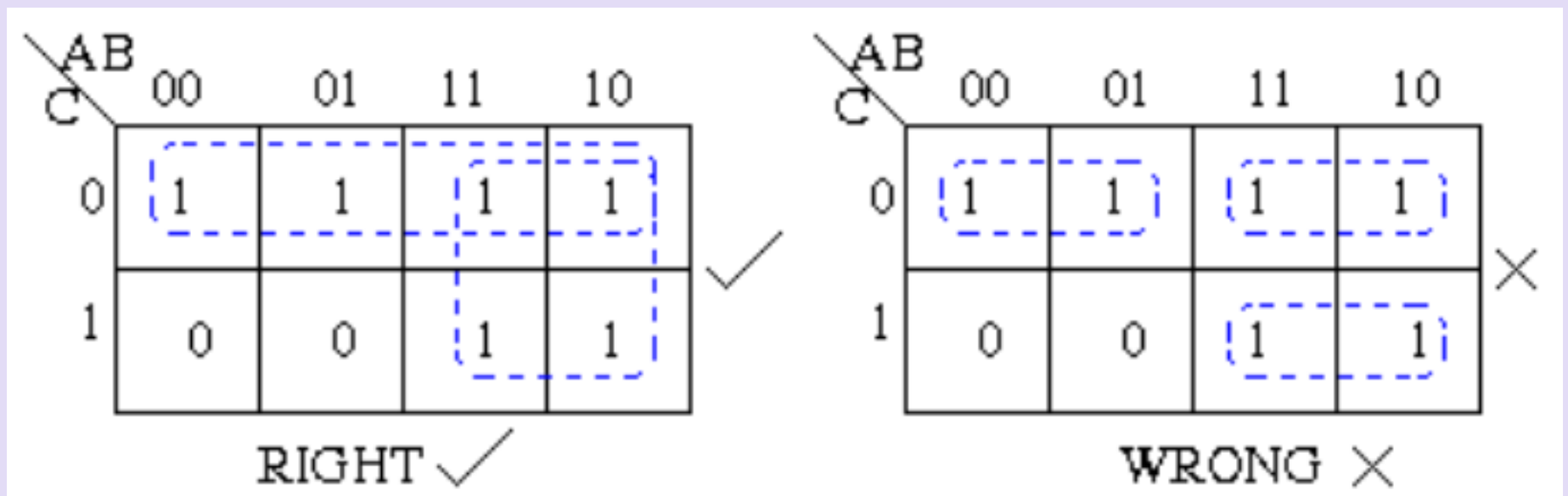
AB C	00	01	11	10	
0	1	1	1	1	← Group of 5
1	0	0	0	1	

WRONG ✗

# Rules for Using K-Maps for Simplification

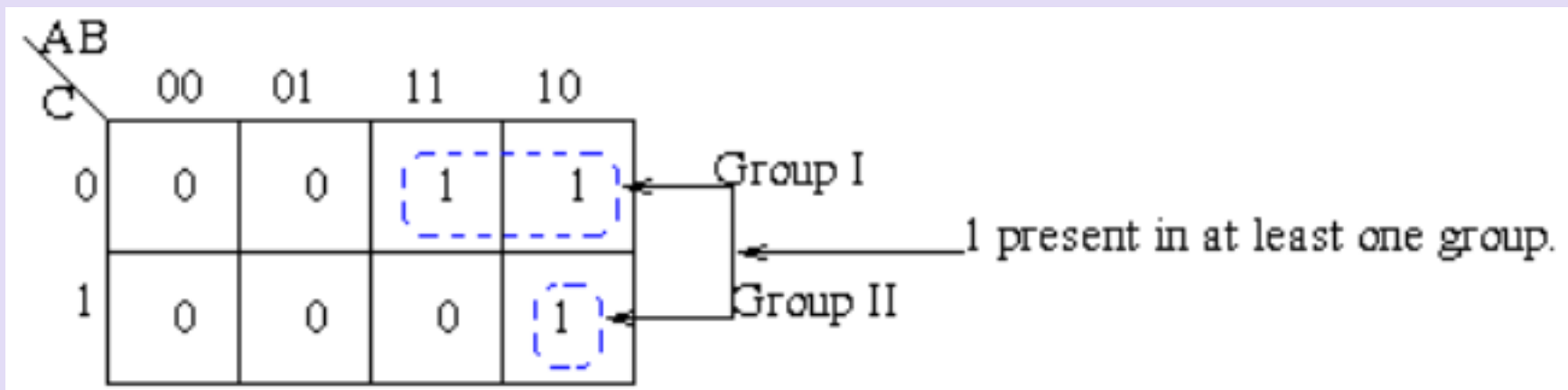
## 5. Each group must be as large as possible

(Otherwise we're not being as minimal as we can be, even though we're not breaking any Boolean rules)



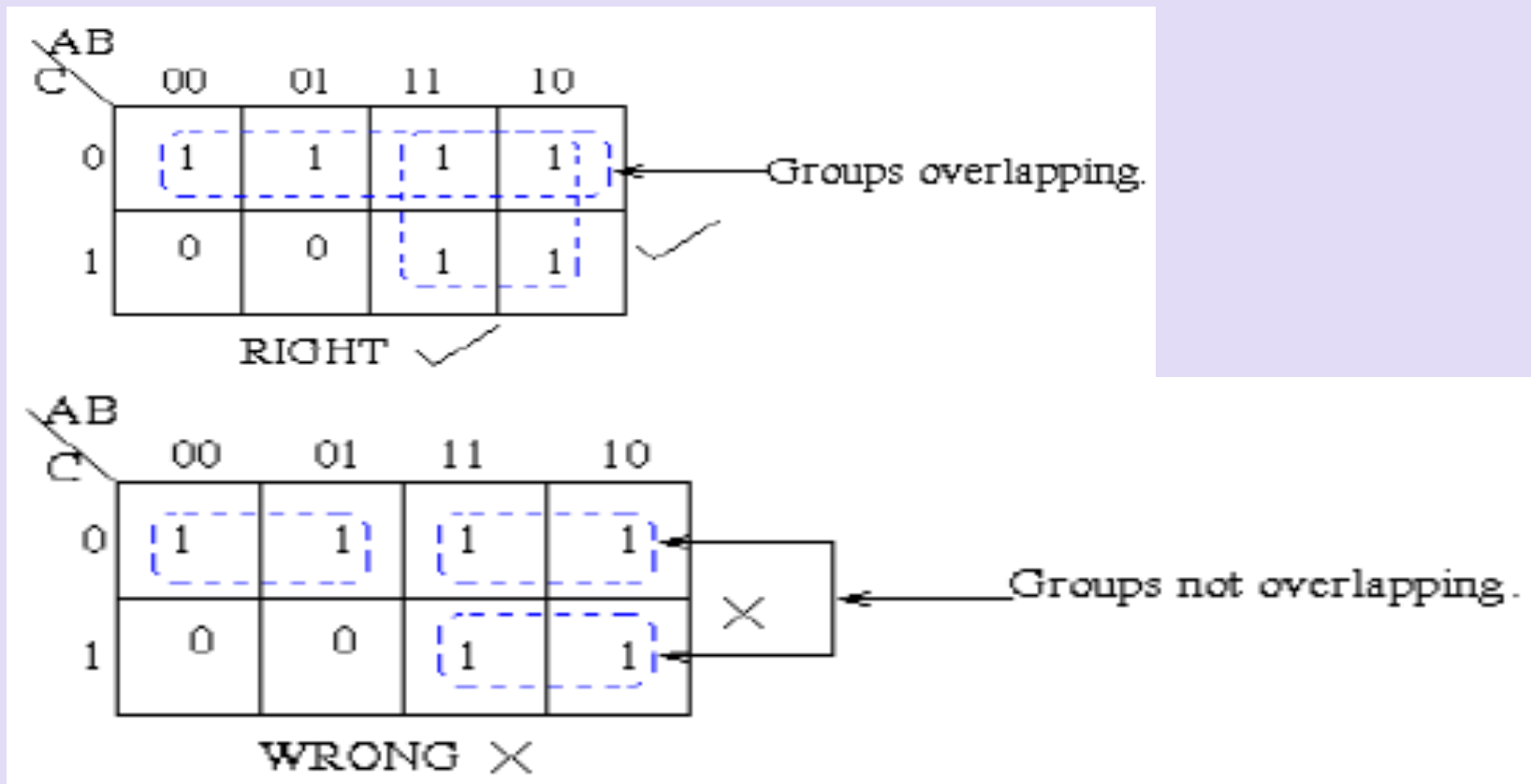
# Rules for Using K-Maps for Simplification

6. Each cell containing a “1” must be at least in one group



# Rules for Using K-Maps for Simplification

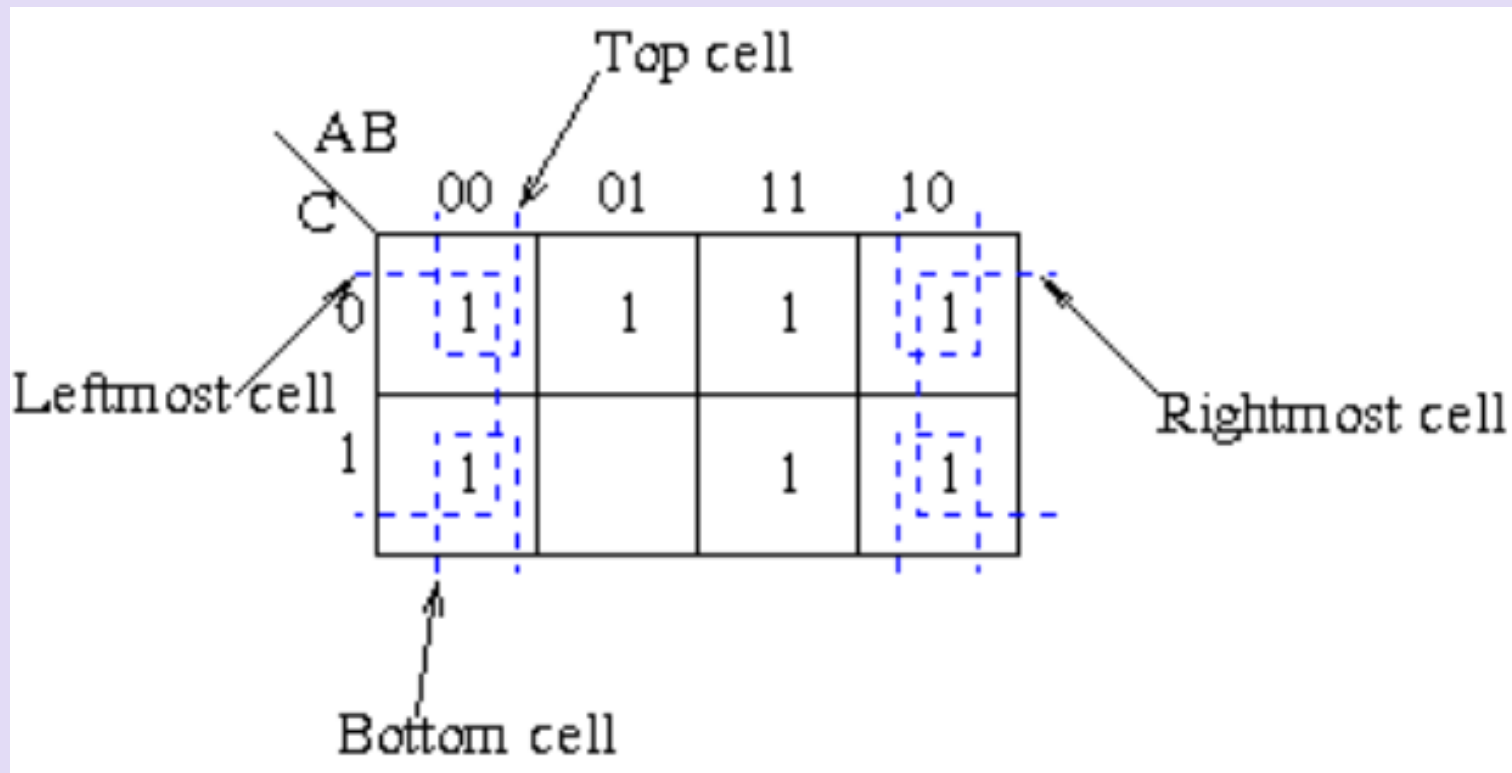
## 7. Groups may overlap esp. to maximize group size



# Rules for Using K-Maps for Simplification

## 8. Groups may wrap around the table.

The leftmost cell in a row may be grouped with the rightmost cell and the top cell in a column may be grouped with the bottom cell.



# Example 1

## 2 vars

$F(X,Y)$

$$= XY + Y$$

$$= Y(X + 1)$$

$$= Y$$

*Verifying results!*

$$F(X,Y) = Y$$

$Y = 1$  column

$X \backslash Y$	0	1
0		1
1		1

# Example 2

## 3 vars

$F(X,Y,Z)$

$$= XZ + Z(X' + XY)$$

$$= XZ + ZX' + ZXY$$

$$= Z(X + X' + XY)$$

$$= Z(1 + XY)$$

$$= Z$$

Verifying results!

$$F(X,Y,Z) = Z$$

A Karnaugh map for the function F(X,Y,Z) with variables X, Y, and Z. The vertical axis is labeled Z, with values 0 and 1. The horizontal axis is labeled XY, with values 00, 01, 11, and 10. The map shows a single group of 1s in the row where Z=1, covering all four XY combinations (00, 01, 11, 10). A red dashed box highlights this group. Above the map, a blue bracket labeled Y=1 spans the 01 and 11 columns, and a yellow bracket labeled X=1 spans the 11 and 10 columns. A red arrow points from the group of 1s to the equation F(X,Y,Z) = Z.

Z \ XY	00	01	11	10
0				
1	1	1	1	1

# Example 3

3 vars

Class Ex.

$$!A!B!C + !A!BC + !ABC + !AB!C + A!B!C + AB!C$$

		AB			
		00	01	11	10
C	0	1	1	1	1
	1	1	1		

$$F(X,Y,Z) = !C + !A$$



# Example 4

## 4 vars

$$F(A,X,Y,Z)$$

$$= AX + Z(X+A'+Y)$$

$$= AX + ZX + ZA' + ZY$$

$$F(A,X,Y,Z) = ZA' + AX + ZY$$

		XY			
		00	01	11	10
AZ	00				
	01	1	1	1	1
	11		1	1	1
	10			1	1

Annotations on the Karnaugh map:
 

- A blue bracket above the 01 and 11 columns is labeled  $Y=1$ .
- A yellow bracket above the 11 and 10 columns is labeled  $X=1$ .
- A blue bracket to the left of the 01 and 11 rows is labeled  $Z=1$ .
- A yellow bracket to the left of the 11 and 10 rows is labeled  $A=1$ .
- Red dashed boxes highlight the 1s in the 01 and 11 rows, the 11 and 10 columns, and the 11 and 10 cells of the 11 row.
- A red arrow points from the  $Z=1$  label to the 11 and 10 cells of the 01 row.

# Example 4

4 vars

Class Ex.

$F(A,B,C,D)$

$$= ABCD' + ABC'D + CD + A'B' + C'D$$

		AB			
		00	01	11	10
CD	00	1			
	01	1	1	1	1
	11	1	1	1	1
	10	1		1	

$B = 1$  (bracket over 01, 11)  
 $A = 1$  (bracket over 11, 10)  
 $D = 1$  (bracket over 01, 11 rows)  
 $C = 1$  (bracket over 11, 10 rows)

$$F(A,B,C,D) = A'B' + D + ABC$$

# K-Map Rules Summary

1. Groups can contain only 1s
2. Only 1s in adjacent groups are allowed
3. Groups may ONLY be horizontal or vertical (no diagonals)
4. The number of 1s in a group must be a power of two (1, 2, 4, 8...)
5. Groups must be as large AND as few in no.s as “legally” possible
6. All 1s must belong to a group, even if it’s a group of one element
7. Overlapping groups are permitted
8. Wrapping around the map is permitted

# Exploiting “Don’t Cares”

- An *output variable* that’s designated “don’t care” (symbol = X) means that it could be a **0** or a **1** (i.e. we “don’t care” which)
  - That is, it is **unspecified**,  
usually because of invalid inputs

# Example of a Don't Care Situation

- Consider coding all decimal digits (say, for a digital clock app):
  - 0 thru 9 --- requires how many bits?
    - 4 bits
  - But! 4 bits convey more numbers than that!
    - Don't forget A thru F!
- Not all binary values map to decimal



## Example Continued...

Binary	Decimal
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7

Binary	Decimal
1000	8
1001	9
1010	X
1011	X
1100	X
1101	X
1110	X
1111	X

# Don't Care: So What?

- Recall that in a K-map, we can only group 1s
- Because the value of a don't care is irrelevant, we can treat it as a 1 ***if it is convenient to do so*** (or a 0 if that would be more convenient)

# Example

- A circuit that calculates if the 4-bit binary coded single digit decimal **input % 2 == 0**
- So, although 4-bits will give me numbers from 0 to 15, I don't care about the ones that yield 10 to 15.

I3	I2	I1	I0	R
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	1
0	1	1	1	0
1	0	0	0	1
1	0	0	1	0
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X



# Example as a K-Map

		$I_1 I_0$			
		00	01	11	10
$I_3 I_2$	00	1	0	0	1
	01	1	0	0	1
	11	X	X	X	X
	10	1	0	X	X

# If We Don't Exploit "Don't Cares"

$$R = \bar{I}_1 \bar{I}_0 \bar{I}_3 + I_1 \bar{I}_0 \bar{I}_3 + \bar{I}_0 \bar{I}_1 \bar{I}_2 I_3$$

		$I_1 I_0$			
		00	01	11	10
$I_3 I_2$	00	1	0	0	1
	01	1	0	0	1
	11	X	X	X	X
	10	1	0	X	X

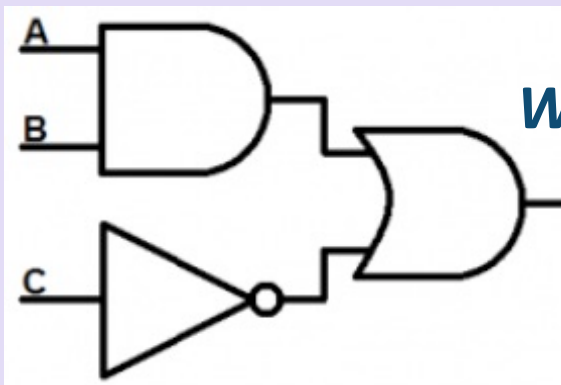
# If We **DO** Exploit “Don’t Cares”

**$R = \bar{I}_0$**

$I_3 I_2$ \ $I_1 I_0$		$I_1 I_0$			
		00	01	11	10
$I_3 I_2$	00	1	0	0	1
	01	1	0	0	1
	11	X	X	X	X
	10	1	0	X	X

# Combinatorial Logic Designs

- When you *combine* multiple logic blocks together to form a more complex logic function/circuit



What is the output?

$$A.B + \overline{C}$$

What is its truth table?

A	B	C	F
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

What is its K-Map?

		00	01	11	10
0		1	1	1	1
1				1	

# Exercise 1

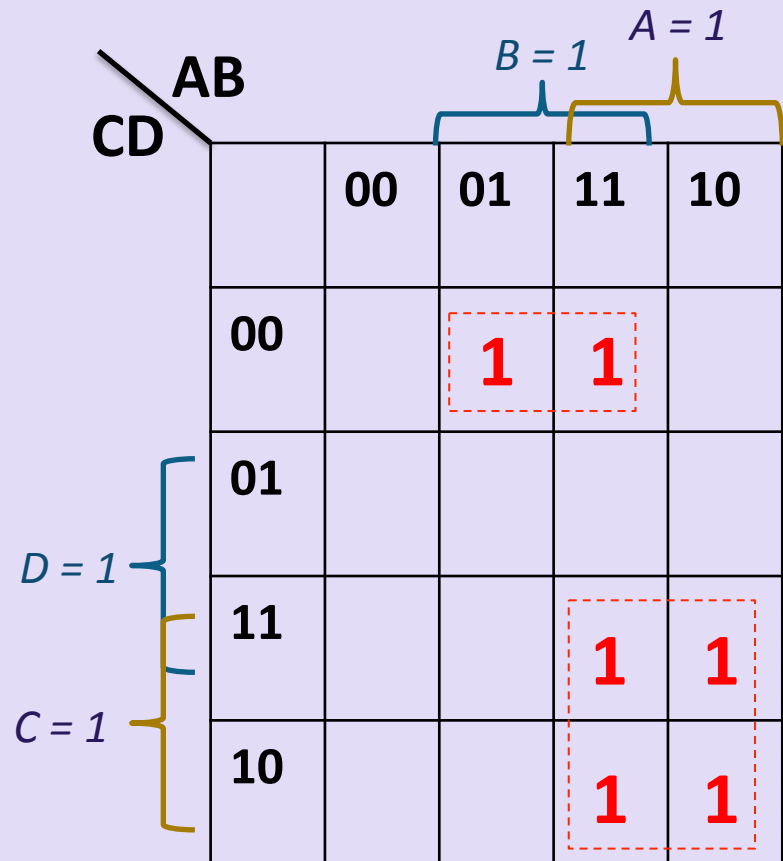
- Given the following truth table, draw the resulting logic circuit
  - **STEP 1:** Draw the K-Map and simplify the function
  - **STEP 2:** Construct the circuit from the now simplified function

A	B	C	D	F
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	1

A	B	C	D	F
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	1

# Exercise 1 – Step 1

*Get the simplified function*

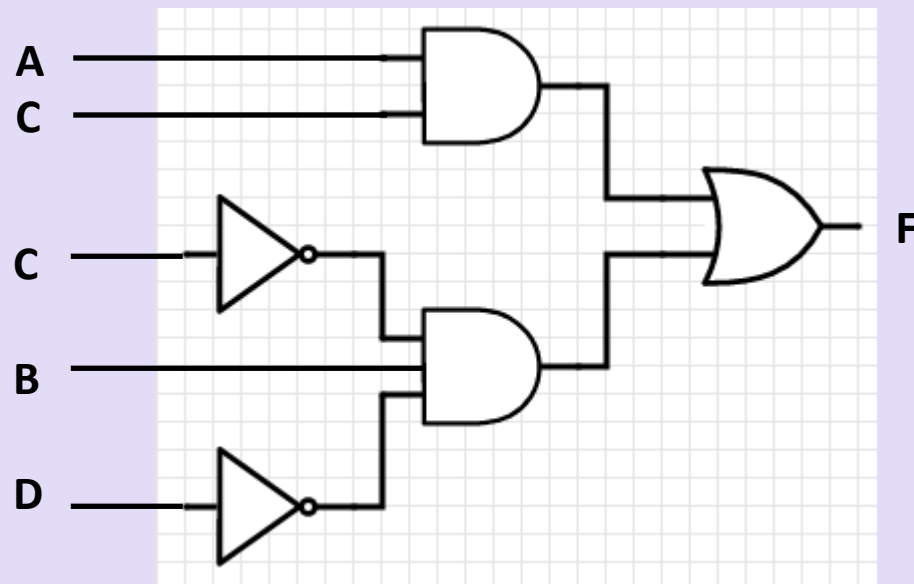


$$F(A,B,C) = B.C'.D' + A.C$$

# Exercise 1 – Step 2

*Draw the logic circuit diagram*

$$F(A,B,C) = B.C'.D' + A.C$$



# Exercise 2

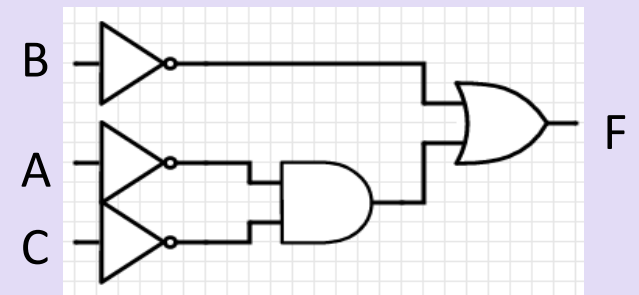
# Class Ex.

- Given the following truth table, draw the resulting logic circuit

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

		AB			
C		00	01	11	10
	0		1	1	
1		1			1

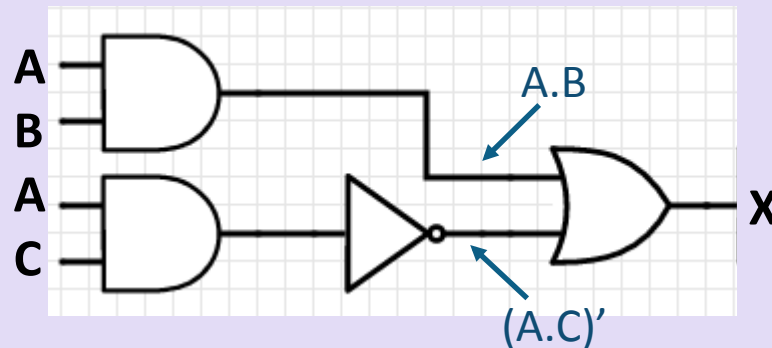
$$F(A,B,C) = B' + A'.C'$$





## Exercise 3

- Given the following schematic of a circuit, (a) write the function and (b) fill out the truth table:



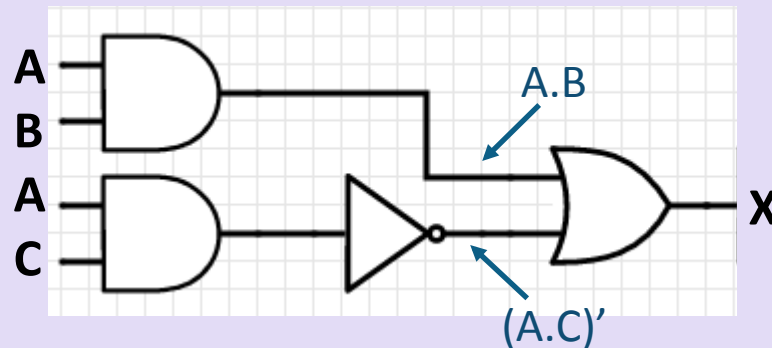
$$X = A.B + (A.C)'$$

(note that also means:  $X = A.B + A' + C'$ )

A	B	C	X
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

## Exercise 3

- Given the following schematic of a circuit, (a) write the function and (b) fill out the truth table:



$$X = A.B + (A.C)'$$

(note that also means:  $X = A.B + A' + C'$ )

A	B	C	X
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

# YOUR TO-DOs

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- Lab 6!

**</LECTURE>**