# Practice Questions (and Answers) for Midterm Exam CS 64, Fall 2019, Matni

<u>IMPORTANT NOTE</u>: These questions are NOT representative of EVERYTHING you need to study for the midterm exam! You should also review your lab assignments questions and also all the examples and demos done in class.

# A. Data Representation

- 1. Convert the following decimal numbers into signed 8-bit binary and two-digit hexadecimal.
  - a. 39
  - b. 104
  - c. -59
  - d. -98
- 2. Convert the following hexadecimal numbers into decimal.
  - a. 0x46
  - b. 0x1D
  - c. 0x3F2
- 3. Convert these unsigned binary numbers into decimal.
  - a. 1010
  - b. 10101010
  - c. 11000001
- 4. Convert these signed binary numbers into decimal.
  - a. 1010
  - b. 00110100
  - c. 11000001
- 5. In 5 bits, what is the most negative value and the most positive value representable in signed form, using two's complement? Express your answers in both binary and decimal.
- 6. Same question as #5, but representable in unsigned form?
- 7. Suppose you are given the following 4-bit binary number, shown in two's complement: **1001**You're not told whether or not the number is signed or unsigned. Is this information important in knowing what the value of the number is, in decimal? That is, do you need to know if it's signed or unsigned to say what the decimal value is? Why or why not?

# **B.** Binary Arithmetic

- 8. Perform the following two's complement addition, noting whether or not the carry bit (C) and/or the overflow bit (V) get set: **01111111 + 11111111**
- 9. Perform the following two's complement addition, noting whether or not the carry bit (C) and/or the overflow bit (V) get set: **00100101** + **10110111**
- 10. Perform the following two's complement addition, noting whether or not the carry bit (C) and/or the overflow bit (V) get set: **11100101** + **10000000**

# C. Bitwise Operations

11. Given the hexadecimals X and Y in each of the following below, find:

```
X \&\& Y, X | | Y, and X \land Y. Give your answers in hexadecimal.
```

```
a. X = 0x2D, Y = 0xFE
```

- b. X = 0x67, Y = 0x67
- c. X = 0x4A, Y = 0x91
- 12. Consider the following C code, which is intended to extract 7 bits from bits position 7 thru 13 of the given input i, leaving the bits in their original position:

```
int unsignedBits7through13(int i) {
  return _____;
}
```

Fill in the blank with a **single** bitwise expression which will make the code do what it is intended to do.

13. Consider the following C code, which is intended to extract the next 7 bits of the given input i, treating the result as a signed value, putting the bits in the rightmost position:

```
int signedBits7through13(int i) {
   // Some lines of code here
   return u;
}
```

Fill in the empty part with valid C/C++ code (can be multiple lines) which will make the code do what it is intended to do.

# D. Assembly

- 14. Why isn't **1i** an actual MIPS instruction?
- 15. Translate the following pseudo-C code into MIPS assembly. Where << read integer from the user>> is used, you should use special functionality provided by SPIM to read in an integer from the console. Where << print integer s1>> is used, you should use special functionality provided by SPIM to print the integer stored in s1 to the console. The variables used below should be placed in the register with the same name. For example, variable s0 should be placed in register \$s0. If you need additional registers than what the code below uses, use registers \$t0 \$t9. You do **not** need to exit the program properly.

```
int s0 = <<read integer from the user>>;
int s1 = 2;
if (s0 < 7) {
   s1 = 3;
} else {
   s1 = s0 + s0;
}
<<pre><<pre><<pre><<pre><<pre>from the user>>;
```

- 16. Write an entire assembly program that takes an integer from standard input (i.e. the user) and checks to see if it is a 0, or 1, or 2 (assume the user behaves and only enters one of these 3 numbers and you do not have to check if they do otherwise). The program then prints to the standard output the text "zero", or "one", or "two", as the appropriate case may be. This has to be followed by a newline character. Bonus points if you can do this using only 2 branch statements.
- 17. What is the machine code (in hexadecimal) for these instructions?

```
a. add $a0, $s3, $a1b. andi $t2, $t2, 119
```

18. Given the C++ function described below:

```
int thatone(int num1, int num2) {
    ans = 4*num1 + 2*num2;
    return (ans);
}
```

- a. Which registers would you use for **num1**, **num2**, and **ans**, assuming MIPS Calling Convention rules?
- b. What assembly instruction would you have at the end of the function?
- c. What assembly instruction would you have in the code where the function is called?

```
19. Consider the C/C++ code below:
    int sum( int n ) {
        if ( n == 0 ) return 0;
        else return (n + sum(n - 1));
    }
```

- a. Knowing that you have to follow the MIPS Calling Convention, which variables should be preserved either directly (via the stack) or indirectly (in an S-register) in order to maintain the intended program behavior?
- b. Implement the previously shown C/C++ code using MIPS assembly, taking care to preserve the values you identified previously. Ignore the .data part and just focus on the .text part of the program. Assume, for this implementation, you have to call the function from your main program like this:

  cout << sum(4)

# **ANSWERS:**

# A. Data Representation

1.

- a. 00100111, 0x27
- b. 01101000, 0x68
- c. 11000101, 0xC5
- d. 10011110, 0x9E

2.

- a. 70
- b. 29
- c. 1010

3.

- a. 10
- b. 170
- c. 193.

4.

- a. -6
- b. 52
- c. -63
- 5. Most negative: 10000, -16. Most positive: 01111, 15.
- 6. You cannot represent negative numbers with unsigned binaries. Most positive: 11111, 31.
- 7. We know it's signed *because it's in two's complement*. If we hadn't been told it was two's complement, then because of the left-most bit being a 1, we would be unable to determine the value of the number unless more information was given. This is because the left-most bit would make it a negative number if the number was signed, and a different positive number if the number was unsigned.

```
B. Binary Arithmetic
```

```
8. 01111111
   +111111111
   -----
    01111110
   The carry bit gets set (i.e. C = 1, V = 0)
9. 00100101
   +10110111
    11011100
   Neither gets set (i.e. C = 0, V = 0)
10. 11100101
   +10000000
   -----
    01100101
   Both get set (i.e. C = 1, V = 1)
C. Bitwise Operations
11.
      a. 0x2C, 0xFF, 0xD3
      b. 0x67, 0x67, 0x00
      c. 0x00, 0xDB, 0xDB
12. (i & 0x00003F80)
13. int signedBits7through13(int i) {
      int u = (i \& 0x00003F80) >> 7;
      if (u & 0x00000040) {
        u = 0xFFFFFF80;
      }
```

# D. Assembly

}

return u;

14. All MIPS instructions are exactly 32 bits large. li can be used to load a 32 bit constant into a register. Therein lies a problem: the constant's entire value couldn't possibly be held in a single instruction, because there aren't enough bits (in addition to the 32 bits of the constant, there would need to be other bits to encode what the instruction is, which requires in total more than 32 bits). It is for this reason that li is a *psuedoinstruction* which can be automatically translated to multiple instructions, whenever we need to specify a constant using the full 32 bits.

```
15.
  main:
   # read in the integer from the user, and initialize s1
     li $v0, 5
     syscall
    move $s0, $v0
    li $s1, 2
    # check if $s0 < 7
     li $t0, 7
     slt $t1, $s0, $t0
    # jump to the else branch if this isn't true
    beq $t1, $zero, else_branch
    # fall through to the true branch
     li $s1, 3
     j print
   else branch:
     add $s1, $s0, $s0
    # fall through to the print
   print:
    li $v0, 1
    move $a0, $s1
     syscall
16.
   .data
   zero: .asciiz "zero\n"
   one: .asciiz "one\n"
  two: .asciiz "two\n"
   .text
  main:
       # Get user standard input
       li $v0, 5
       syscall
      move $s0, $v0
       li $t0, 0
       li $t1, 1
       li $t2, 2
       # Assume that input is 0
       # Check to see if input is 1 or 2 and branch accordingly
       beq $s0, $t1, print1
      beq $s0, $t2, print2
```

c. jal thatone

```
# Print "zero"
       li $v0, 4
       la $a0, zero
       syscall
       j end
   print1: # Print "one"
       li $v0, 4
       la $a0, one
       syscall
       j end
   print2: # Print "two"
       li $v0, 4
       la $a0, two
       syscall
   end:
       li $v0, 10
       syscall
17.
      a. 0x02652020
      b. 0x314A0077
18.
      a. $a0 for num1, $a1 for num2, $v0 for ans
      b. jr $ra
```

```
19.
      a. a0 for n, v0 for the returned value (n + sum(n - 1))
                                                    int sum( int n ) {
      b.
                                                       if ( n == 0 ) return 0;
                                                       else return (n + sum(n - 1));
         .text
      sum:
          addiu $sp, $sp, -8
                                   # PUSH
          sw $ra, 4($sp)
          sw $s0, 0($sp)
          beq $a0, $zero, return # is size !=0?
                               # preserve a0 (variable n)
          move $s0, $a0
                               # add n to $v0
          add $v0, $v0, $s0
          addi $a0, $a0, -1
                             # n is now: n - 1
          jal sum
                               # recursive call
      return:
          lw $ra, 4($sp)
                                   # POP
          lw $s0, 0($sp)
          addiu $sp, $sp, 8
          jr $ra
      main:
         li $v0, 0
                                   # Initialize sum ($v0)
         li $a0, 4
                                   \# n = 4
         jal sum
                                   # Call sum(4); expect $v0 to be 10
         move $a0, $v0
         li $v0, 1
         syscall
         li $v0, 10
         syscall
```